



Kryptoagilität in der Standardisierung der IETF

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Die neuen PQC Algorithmen

Das OpenPGP Protokoll

PQC für OpenPGP

PQC für X.509 und CMS

Agilität bei Sicherheitseigenschaften bei OpenPGP und X.509 / CMS

OpenPGP und LibrePGP

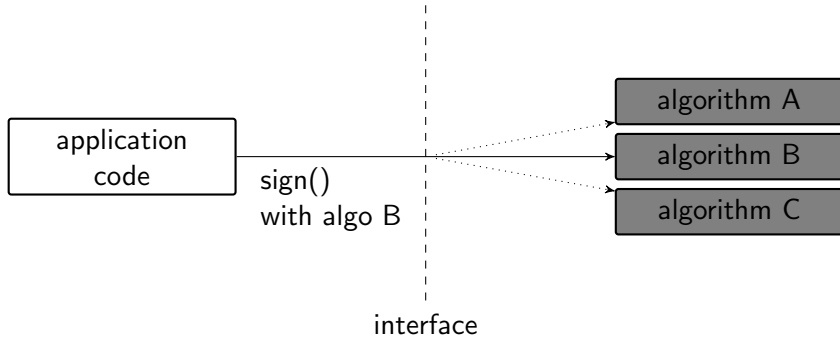
Zusammenfassung

The new PQC Algorithms

- ▶  KEM (Verschlüsselung)
 - ▶ ML-KEM 
- ▶  Signaturen
 - ▶ ML-DSA 
 - ▶ SLH-DSA 
 - ▶ XMSS / LMS  

Cryptographic Agility relies on stable interfaces

- ▶ To easily exchange an algorithm, the interface must not change
- ▶ Constant interface definition: treat algorithm as black box



Price question

Which interface gets “more complicated” through PQC? Encryption or signatures?

KEM Interfaces

► KEM

- Interface change compared to encryption scheme
- KEM encrypts a random key
- random key used to encrypt session key or data directly

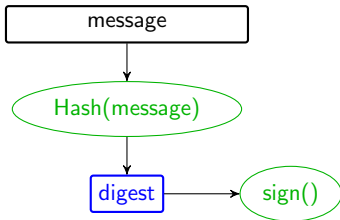


Complexity of KEM interface

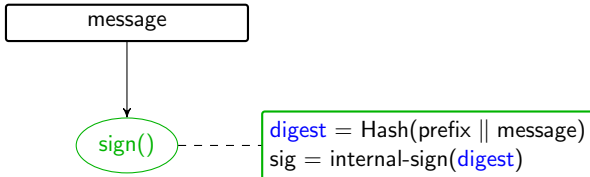
- ▶ currently only one KEM: ML-KEM
- ▶ no interface changes to be expected for future KEMs

Signatures: From “hash-then-sign” to “pure signing”

- ▶ traditional algorithms RSA, ECDSA:
 - ▶ build by Hash-then-sign paradigm
 1. hash the message stream
 2. sign the hash



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- ▶ modern / PQC design paradigm is not hash-then-sign
 - ▶ advantage: increased collision resistance of signature scheme
 - ▶ PQC signature algorithms receive the message stream directly



Signature Interface

	NIST FIPS std.	IETF (RFCs)
schemes	ML-DSA, SLH-DSA	XMSS, LMS
context parameter	✓	✗
pre-hash variant ¹	✓	✗

¹pre-hash variant $\approx$ hash-then-sign

Signatures: Online Operations

- ▶ Online property (a.k.a stream processing, incremental processing):
 - ▶ process the message during signing or verification as a stream, requires only **constant** (small) amount of **memory**
- ▶ if not possible, need to store full message
 - ▶ sometimes impossible for resource-constrained devices

- ▶ $\mu \leftarrow H(H(\text{public-key}) \parallel M)$ // *hedges against multi-user attacks*
- ▶ then sign μ
- ▶ short message representative μ
- ▶ FIPS: computing μ needs to be done in a cryptographic module

can compute message hash as a stream

- 
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 - ▶ then sign μ
 - ▶ short message representative μ
 - ▶ FIPS: computing μ needs to be done in a cryptographic module

SLH-DSA

- ▶ $R \leftarrow \text{PRF}_{\text{msg}}(SK.\text{prf}, \text{opt_rand}, \text{PK.root}, M)$ // *M to hedge against RNG failure*
- ▶ $\text{digest} \leftarrow H_{\text{msg}}(R, \text{PK.seed}, \text{PK.root}, M)$ // *improved collision resistance*
- ▶ compute signature based on *digest* ...

SLH-DSA

cannot process message hash as a stream;
need to store M for second computation

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SLH-DSA

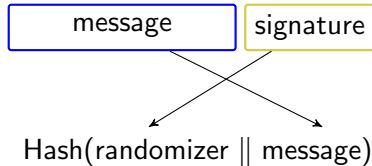
cannot process message hash as a stream;
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- ▶ $R \leftarrow \text{PRF}_{\text{msg}}(\text{SK.prf}, \text{opt_rand}, \text{PK.root}, M)$ // M to hedge against RNG failure
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- ▶ compute signature based on *digest* ...

- pure SLH-DSA: signing, verification not online, no streaming
- need to resort to pre-hash (hash-then-sign) variant for memory efficiency
 - no increased collision resistance at all

Hash-based signature verification

- ▶ all hash-based signature schemes prefix a randomization string (“randomizer”) to the signed data
- ▶ randomizer is encoded in the signature



Properties of signatures schemes

	Algorithm specification (NIST, RFC) (final)				📄 X.509 (cert, CRL) specification (LAMPS drafts, not final)			
	ext. msg. repr.	pre- hash	sign- online	verify- online	ext. msg. repr.	pre- hash	sign- online	verify- online
ML-DSA	✓ ^(μ)	✓	✓ ²	✓ ²	✓ ^(μ)	✗	✓ ²	✓ ²
SLH-DSA	✗	✓	(✓) ³	(✓) ⁴	✗	✓	(✓) ³	(✓) ⁴
XMSS/LMS	✗	✗	✓	✗ ⁵	✗	✗	✓	✗

² initialisation of ML-DSA stream-processing for message requires public-key

³ online-signing with SLH-DSA only with pre-hash variant

⁴ online-verification with SLH-DSA only with pre-hash (or signature-first or randomizer-first)

⁵ XMSS/LMS online-verification only with pre-hashing on application level or with signature-first

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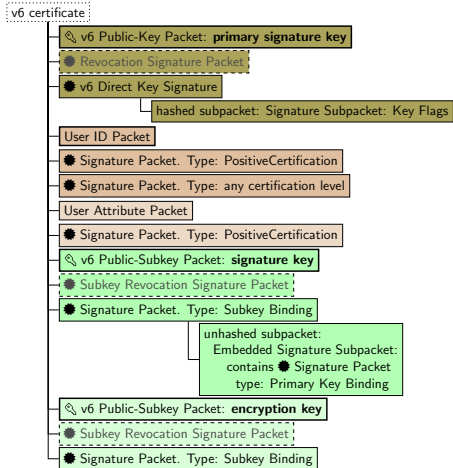
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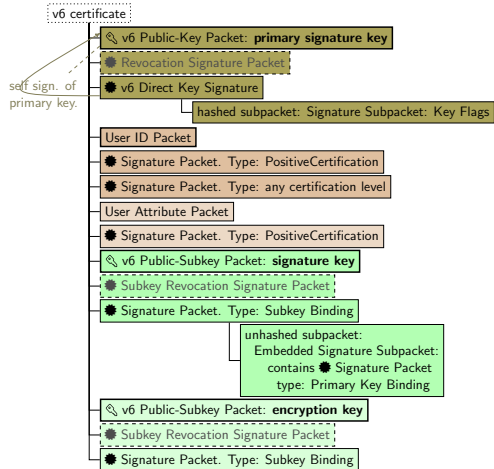
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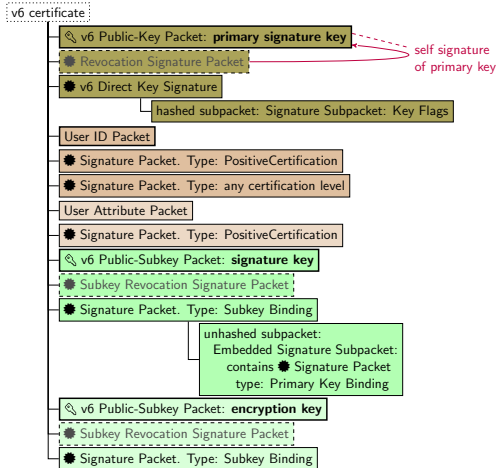
Zusammenfassung

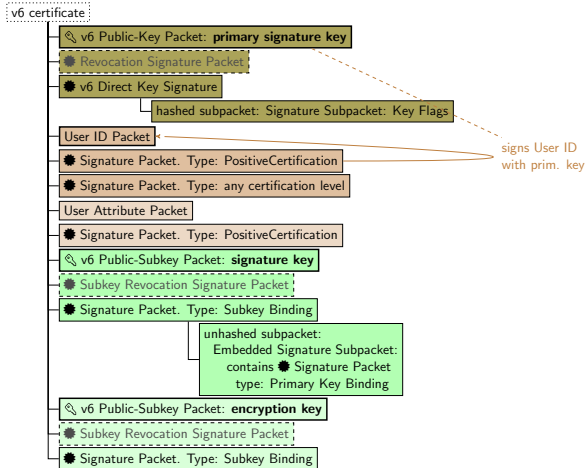
OpenPGP Protocol

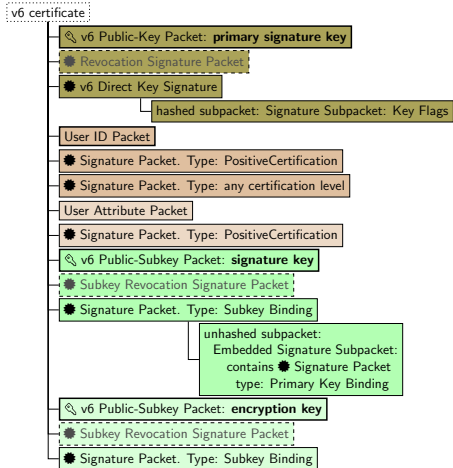


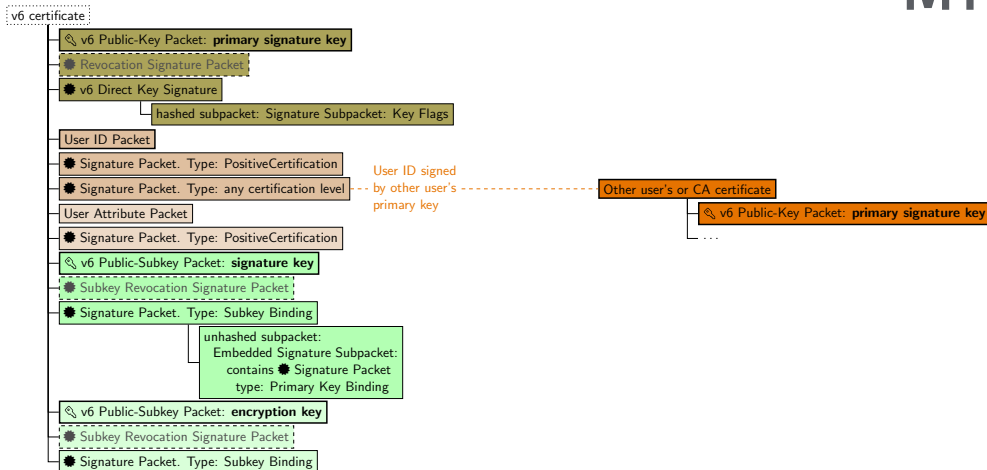


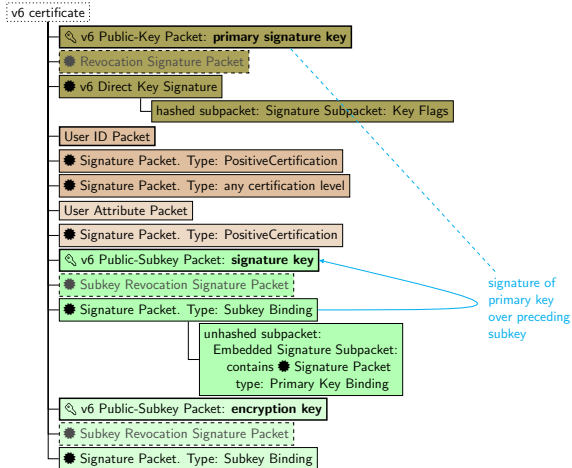


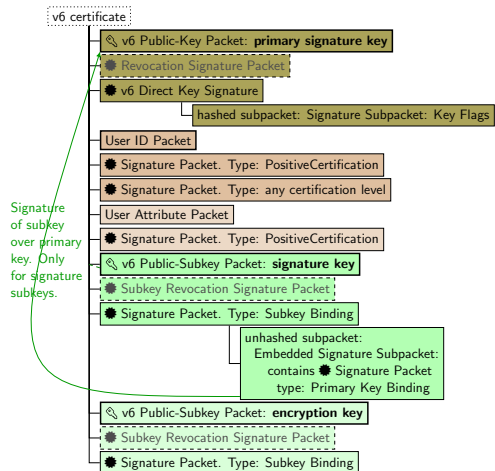


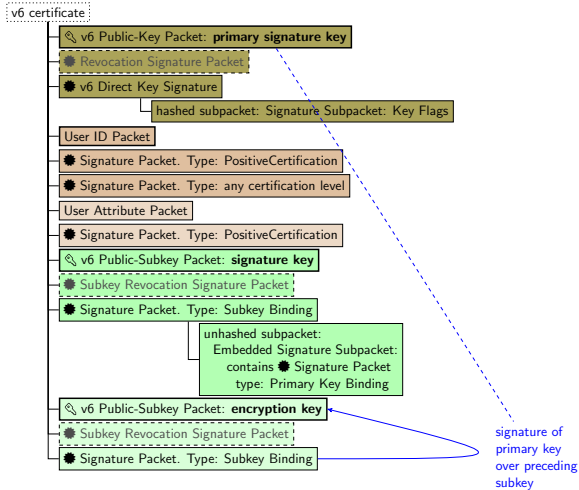


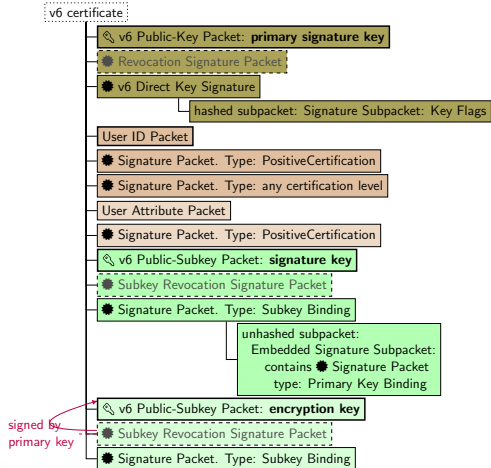


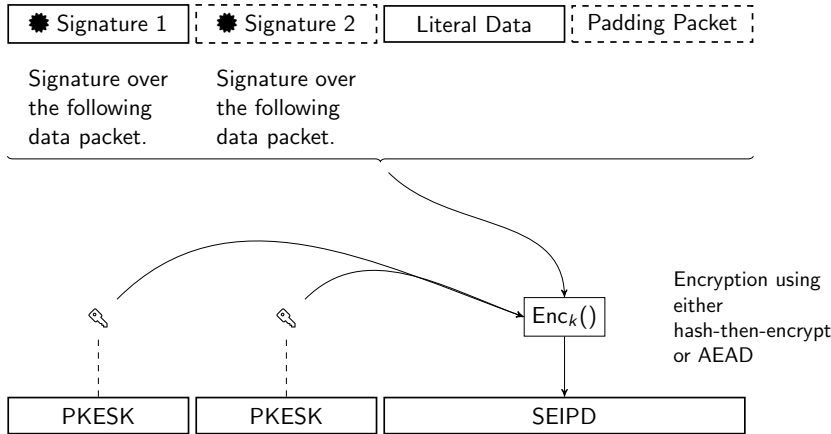












Public Key Encrypted Session Key.

session key encrypted to public key of user A

session key encrypted to public key of user B

Symmetrically Encrypted Integrity Protected Data Packet.

(message encrypted with session key)

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PQC@Thunderbird

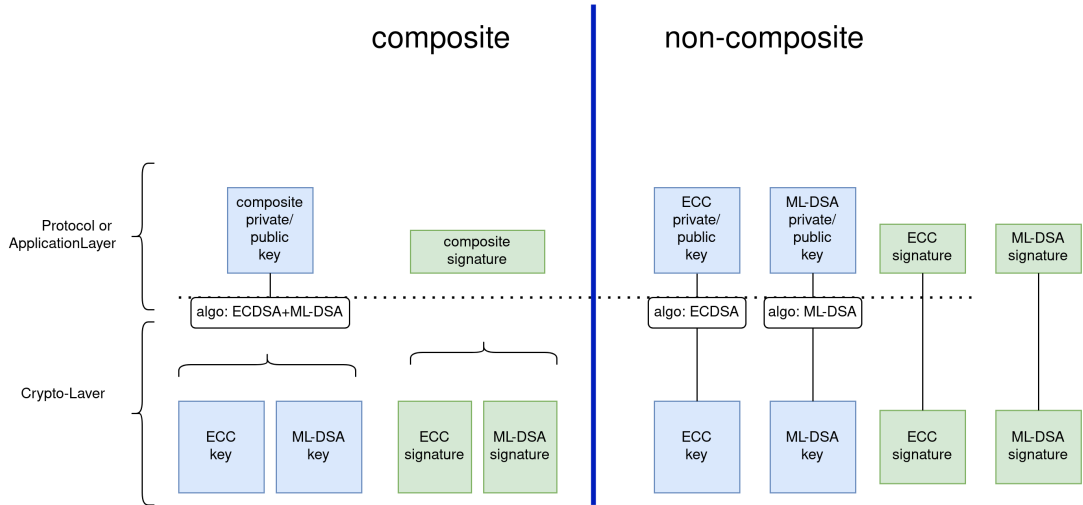
- ▶ BSI Projekt 480 “PQC@Thunderbird”
- ▶ Standardisierung von PQC in OpenPGP
- ▶ Implementierung in Thunderbird
- ▶ MTG mit TU Eindhoven
- ▶ Dez. 2021 - Nov. 2025

PQ/T hybrid schemes

- ▶ European governments recommend/require pairing PQC schemes with traditional algorithms
 - ▶ Exception: Hash-based algorithms **#**
- ▶ NIST: “new stuff sometimes gets broken”⁶

⁶<https://csrc.nist.gov/csrc/media/Presentations/2025/draft-sp-800-227-overview/images-media/sp-800-227-galagic.pdf#page=11>

PQ/T hybrid signatures



PQC-Integration: Ursprüngliche Idee aus dem Projekt

- ▶ in jedem Fall: hybrid = “multi-algorithm” (außer SLH-DNA)
- ▶ erster Ansatz: vollständig generisch
 - ▶ Neue Algorithmen ID für ML-KEM, ML-DNA, SLH-DNA
 - ▶ Non-composite für generische Kombinationen von Algorithmen
 - ▶ Mehrere Signaturen in OpenPGP schon vorgesehen
 - ▶ z.B. RSA + ML-DNA unter Angabe der jeweiligen Algorithm IDs
 - ▶ und unter Angabe der Parameter:
 - ▶ RSA **4096**
 - ▶ ML-DNA **192 bit**

Umentscheidungen bei PQC Integration

- ▶ Umentscheidung 1
 - ▶ Composite
 - ▶ Algorithm ID: Feste Kombinationen aus PQ / T
- ▶ Umentscheidung 2
 - ▶ Sicherheitsparameter auch mit Algorithm ID festschreiben
 - ▶ Algorithm ID = 30 → “ML-DSA-65+Ed25519”

Two PQC Drafts

 draft-ietf-openpgp-pqc

(adopted, WGLC)

BSI, MTG, Proton

algorithms	requ.	security
ML-KEM-768+X25519	MUST	192 / 128
ML-KEM-1024+X448	SHOULD	256 / 224
ML-DSA-65+Ed25519	MUST	192 / 128
ML-DSA-87+Ed448	SHOULD	256 / 224
SLH-DSA-SHAKE-128s	MAY	128
SLH-DSA-SHAKE-128f	MAY	128
SLH-DSA-SHAKE-256s	MAY	256

 draft-ehlen-openpgp-nist-bp-comp

(not adopted)

BSI, MTG, NIST

algorithms	security
ML-KEM-512+ECDH- NIST-P-256	128 / 128
ML-KEM-768+ECDH- NIST-P-384	192 / 192
ML-KEM-1024+ECDH- NIST-P-384	256 / 192
ML-KEM-768+ECDH- brpP256r1	192 / 128
ML-KEM-1024+ECDH- brpP384r1	256 / 192
ML-DSA-44+ECDSA- NIST-P-256	128 / 128
ML-DSA-65+ECDSA- NIST-P-384	192 / 192
ML-DSA-87+ECDSA- NIST-P-384	256 / 192
ML-DSA-65+ECDSA- brpP256r1	192 / 128
ML-DSA-87+ECDSA- brpP384r1	256 / 192
all "MAY"	

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Zusammenfassung

LAMPS WG: X.509, CMS, etc.

- ▶ Ähnlichkeiten:
 - ▶ Auch nur feste Algorithmenkombinationen
 - ▶ Auch Festlegung der Sicherheitsparameter durch Algorithm-ID
- ▶ Unterschiede:
 - ▶ XMSS und LMS
 - ▶ Alle 12 SLH-DSA Parameter
 - ▶ ML-DSA und ML-KEM als Standalone
 - ▶ Composite
 - ▶ Kombinationen mit RSA und ECDSA (NIST, Brainpool)
 - ▶ Berücksichtigung des Falcon Signaturverfahrens (noch nicht finalisiert von NIST)

Kryptoagilität in OpenPGP und X.509/CMS

- ▶ “niemand (in der IETF) will freie Kombinierbarkeit der Algorithmen”
 - ▶ composite multi-algorithm
- ▶ OpenPGP: starke Einschränkung der neuen Algorithmen gefordert
 - ▶ Interoperabilität steht im Vordergrund
- ▶ X.509/CMS: Größere Kryptoagilität (?) durch größere Algorithmen- und Parameterauswahl
 - ▶ Große Zahl von abgeleiteten Protokollen
 - ▶ Vielfalt steht im Vordergrund

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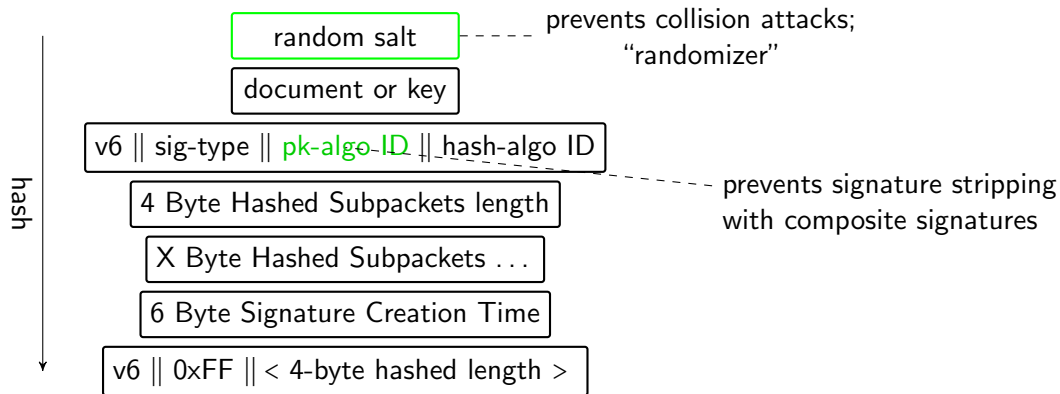
Zusammenfassung

Protocol security features and agility

- ▶ Some protocol security features ...
 - ▶ generally relevant
 - ▶ relevant to PQC
- ▶ How to achieve security updates to protocols?
 - ▶ “the agile way” (?): pick new features explicitly
 - ▶ protocol version upgrade

RFC 9580 OpenPGP – signatures

v6 signatures — hashed data



OpenPGP uses hash-then-sign

- ▶ PQC algorithms used with hash-then-sign
- ▶ no disadvantage due to random salt (= randomizer)
- ▶ good cryptographic agility compared to X.509

Non-Separability in OpenPGP Signatures



v6 signature packet

v6 = 0x06

sig-type 0x00

1B pk-algo = hybrid-...

1B hash-algo

2B hashed subpacket length

hashed subpackets

2B unhashed subpacket length

unhashed subpackets

2B checksum for hash-value

algorithm-specific signature data

sig 1

sig 2

Non-Separability in OpenPGP Signatures



*hashed and
signed
as meta
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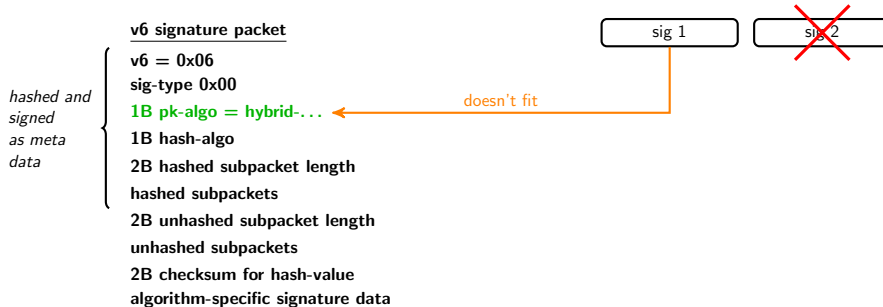
2B checksum for hash-value

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sig 1

~~sig 2~~

Non-Separability in OpenPGP Signatures



CMS – Signatures

- ▶ Cryptographic Message Syntax (e.g. S/MIME)
 - ▶  X.509 certificates
 - ▶  Signatures
 - ▶  Public-key encryption
- ▶ No random salt
 - ▶  Pre-hash signatures potentially vulnerable to hash-collision attacks
- ▶ No meta-data is hashed
 - ▶  Public-key algorithm not fixed by signature
 - ▶ → In case of hybrid: signature stripping attacks require extra countermeasure
 - ▶ Not fully sound if legacy signature remains
- ▶ CMS legacy problem
 - ▶  (theoretical) signature forgery problem with Signed Attributes⁷
 - ▶ Alternative views of what was exactly signed
 - ▶ Was not solved with context parameter of new PQC algorithms

⁷An EUF-CMA violation, see <https://eprint.iacr.org/2023/1801>

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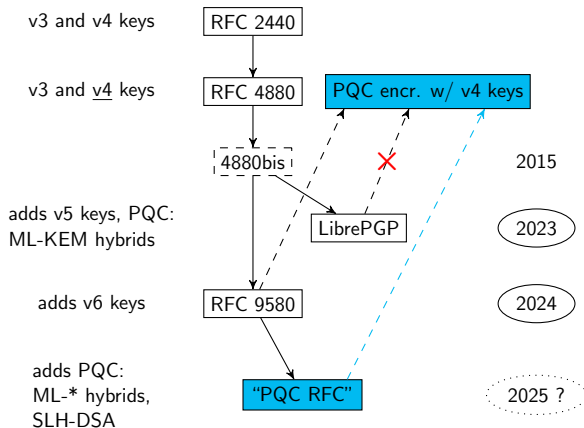
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Zusammenfassung

v6/v5/v4 PQC



- ▶ Main goal: fast adoption of PQC encryption
- ▶ GnuPG LibrePGP standard features incompatible ML-KEM hybrids
 - ▶ LibrePGP (outside IETF) is a recent fork of OpenPGP due to WG-internal unresolved technical "issues"

LibrePGP's approach to PQC

- ▶ So far, only ML-KEM is specified (no signatures)
- ▶ ML-KEM in combinations with ECDH
 - ▶ with any EC curve out 16
- ▶ greater parameter choice than
 - ▶ OpenPGP
 - ▶ X.509 / CMS

Protocol security features in comparison

	CMS	OpenPGP RFC 4880 v4	OpenPGP RFC 9580 v6	LibrePGP
AEAD	⚠ requires opt. mech- anism from RFC 9709 ⁸	⚠ no AEAD	✓ sound AEAD	⚠ un- mitigated weakness ⁸
enhanced signature collision resistance	⚠ not for pre-hash vari- ants or CMS with Signed Attributes ⁹	⚠ no	✓ random salt	⚠ no

⁸<https://cic.iacr.org/p/2/1/19>

⁹Signed Attributes $\approx$ pre-hash

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- ▶ Die neuen PQC Signaturalgorithmen haben Variabilität im Interface
 - ▶ Betrifft v.a. die Speichereffizienz → hauptsächlich relevant für eingebettete Systeme
 - ▶ Hat Einfluss auf die Kryptoagilität (Austauschbarkeit)
 - ▶ Wechsel auf pre-hash kann notwendig werden
- ▶ IETF eher kritisch gegenüber generischer Kombinierbarkeit
 - ▶ bisher (fast) nur Vorschläge für composite multi-algorithm
 - ▶ OpenPGP: vergleichsweise restriktive Algorithmenauswahl, wenige Kombinationen
 - ▶ X.509 / CMS: deutlich größere Auswahl an Algorithmen und Kombinationen
- ▶ Protokolldesign und PQC
 - ▶ OpenPGP: hashing von Metadaten (salt, Algorithm ID)
 - ▶ Erlaubt u.a. Verwendung von hash-then-sign ohne Nachteile
 - ▶ X.509 / CMS
 - ▶ hash-then-sign nur mit schlechterer Sicherheitsaussage
- ▶ Sicherheitsfeatures in Protokollen
 - ▶ Zuverlässige Verbesserungen erfordern Sprung in der Protokollversion
 - ▶ “Agiles Nachrüsten” von Gegenmaßnahme eher problematisch

Thank you for your attention

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